T-DPLL: Online Schema

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1.    SatValue T-DPLL (T-formula φ, T-assignment & μ) {
2.       if (T-preprocess(φ,μ) == Conflict);
3.          return Unsat;
4.       φ^P = T2B(φ); μ^P = T2B(μ);
5.       while (1) {
6.          T-decide_next_branch(φ^P,μ^P);
7.          while (1) {
8.             status = T-deduce(φ^P,μ^P);
9.             if (status == Sat) {
10.               μ = B2T(μ^P);
11.               return Sat; }
12.            else if (status == Conflict) {
13.               blevel = T-analyze_conflict(φ^P,μ^P);
14.               if (blevel == 0)
15.                  return Unsat;
16.               else T-backtrack(blevel,φ^P,μ^P);
17.            }
18.            else break;
19.   } } }
```

Input:

Original T-formula ϕ

Reference to a T-Assignment μ (a set of T-literals) which is initially empty

T-preprocess:

Simplifies ϕ into a simpler version (ϕ^P) using

- Boolean preprocessing from DPLL
- Theory-dependent rewriting steps on T-literals of ϕ

Updates μ if necessary

Returns UNSAT if a conflict is encountered

T-decide-next-branch

Selects literal to branch on as in DPLL

T-deduce

Iteratively deduces the implied Boolean literals and updates ϕ^P and μ^P until:

1. μ^P violates ϕ^P (and entails { [] }) – return Conflict
2. μ^P satisfies ϕ^P, and the result is run through the T-solver
   - satisfied by T-solver? return Sat
   - unsatisfied by T-solver? return Conflict
3. Nothing new can be deduced – return Unknown

T-analyze-conflict

Determine the Boolean conflict set, n^P; add it to ϕ^P, then backtrack up to blevel